Math Logic: Model Theory & Computability

Lecture 01

Model Theory

Mathematical structures.

Every methenetician recognizes a methemetical structure then he the sees Examples. (a) A graph is a pair G := (V, E), there V is a set (of certices) and E = V² is a binary relation on V, called the set of edges $\begin{array}{c|c} \hline 0 & \hline 0 \\ \hline 0 & \hline 0 \\ \hline 2 & \hline 0 \\ \hline \end{array} \end{array} \qquad \begin{array}{c|c} \hline 0 & \hline 0 &$ (b) A partial order is a pair (X, ≤), where X is a set and ≤ is a binary relation on X (i.e. a subset of X²) substying certain axions:
(i) Reflexivity: x ≤ x for all x ∈ X.
(ii) Anti-symmetry: x ≤ y and y ≤ x then x = y for all x, y ∈ X.
(iii) Transitivity: x ≤ y and y ≤ z then x ≤ z for all x, y, z ∈ X. For example the usual \leq on \mathbb{Z} , or the subset relation \subseteq or $\mathbb{P}(\mathbb{N}) :=$ the powerset of a set \mathbb{N} . (c) A group is a quadruple [== (Γ, 1, ., ()⁻¹), where Γ is a set, 1 is an element of Γ, . is a binary operation on Γ, i.e. .: Γ²→Γ, and ()⁻¹ is a unary operation on Γ, i.e. ()⁻¹: Γ→Γ, satisfying the following conditions:

(i) is associative: x. (y. Z) = (x.y). Z for all x, y, ZE [. (ii) 1 is an identity for \bullet : 1 - x = x - 1 = x for all $x \in \Gamma$ (iii) ()^T is an inverse b(\bullet : $x^{-1} - x = x - x^{-1} = 1$ for all $x \in \Gamma$. For example, Z := (Z, 0, +, -()), Sym(n) is a non-abelian group of all permitations of {0, 1, ..., n-1} with the operation of wapori-tion, GL, (IR) := the group of invertible non-matrices over IR under multiplication. (d) A ring is a 6-tuple R:= (R, 0, 1, +, -(), .) where (R, 0, +, -(1)) is an abelial ycoup, 1 ∈ R, . is a binary operation satisfying: (i) . is associative [ii] $1 \text{ is an iduality for }, \text{ i.e. } 1 \cdot x = x \cdot 1 = x \text{ for all } x \in R.$ [iii] $0 \text{ distributes over } +, \text{ i.e. } x \cdot (y + z) = xy + xz \text{ and } (y + z) \cdot x = yx + z \cdot x$ [iv] $0 \neq 1$. for all $x_1 y_1 z \in R$. For excepte, the city of indegers $Z := (Z, 0, 1, t, -(), \cdot)$, the ring of them matrices $M_n(IR)$ over IR, the ring of polyno-mials IF[t] over a field IF, the ring of confirmous tructions f: [0,1] > IR unles usual addition and multiplication, the ring of linear transformations V-SV for a reator space V over a field IF unles the operations of addition (as +) and co-position

(e) A field is a ring E = (IF, 0, 1, +, -(), ·) such that
(i) every non-zero element has a ·· invirse, i.e. for all nonzero x EF
there is g EIF such that x-y=g·x=1.
(ii) · is commutative: x-y= y-x for all x, y E IF.

(as ·).

(f) An ordered field is a 7-taple R := (R, 0, (, t, -(), ·, ≤), where (R, 0, |, t, -(1, ·) is a field, (R, ≤) is a total / hinen order (i.e. a partial order s.t. for all x, y∈ R, x ≤ y or y ≤ x) satisfying: (i) x ≤ y => Z+x ≤ Z+y for cl ×, y, Z ∈ R, (ii) x ≤ y => Z-x ≤ Z-y for cl ×, y, Z ∈ R Ame ZZO. For example the ordered field IR of necls. So it looks like, we can offically define a mathematical structure as a quadruple S == (S, E, E, D), where S is a set, C is a set of constructs from S (possibly empty), & is a (possibly empty) set of operations on S (possibly of different arity = quiphingulunhimpinh), and R is a set of relations on S (again possibly empty). This is an automated definition bease it doesn't provide names for the constants, operations, and relations, head makes it inconvenient to

write axiomic that we want them to satisfy. Thus, we first introduce a naming system/format, and then define structures in a given format.

Dot. A signature or language is a gardaphe $\sigma := (\mathcal{L}, \sigma, \sigma, \sigma)$, there $\mathcal{L}, \sigma, \sigma, \sigma$ are sets of symbols (i.e. meaningless characters, e.g. \mathcal{S}) or names, and $a : \mathcal{FUR} \to \mathbb{N}^{+} := \{1, 2, 3, ...\}$ called arity $(= \{m, m, m, m, m, m, m\})$. We call $\mathcal{L}, \sigma, \sigma$ the sets of contract symbols, function symbols, and celebion symbols. These might be empty.